

Political Economy  
Theory and Experiments  
Lecture 4

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## Information Aggregation

- . One typical feature of democracies is voting.
- . Each citizen out of a large group is asked for her opinion.
- . Opinions are counted and each weighs the same.
- . Is democratic voting better than letting a single competent expert decide (enlightened autocracy)?
- . What are the information aggregation properties of voting?
- . Condorcet supposed that voting will perform better than expert decision, as long as the number of voters is sufficiently large.
- . This, even if every voter has low competence/information.

## Condorcet Jury Theorem

### **The model**

- . There are 2 alternatives  $d \in \{0, 1\}$ . One of them is “right”.
- . The decision is made by majority vote.
- . Voters vote independently of each other, no abstention.
- . The a-priori probability (i.e. before voters get further information) of being right is the same for both alternatives.
- . Each voter has the same probability  $p$  to vote “correctly.”
- . There is an odd number of voters  $n = 2m + 1$ , with  $n \geq 3$ .

## Analysis

. The probability that in the case of  $n$  voters, exactly  $x$  of them make the right decision, depends on  $p$  and is given by:

$$b_n(x; p) = \binom{n}{x} p^x (1-p)^{n-x}$$

with  $x = 0, 1, \dots, n$  and  $p \in (0, 1)$ .

. The probability a majority vote selects the right decision is:

$$M_{2m+1}(p) \equiv \Pr(X \geq m+1) = \sum_{x=m+1}^{2m+1} b_{2m+1}(x; p)$$

. This function is symmetric:  $M_{2m+1}(p) = 1 - M_{2m+1}(1-p)$ .

. Thus, we have:  $M_{2m+1}(1/2) = 1/2$ .

- . How does probability to make right decision by majority voting change with more voters?
- . Suppose we add two voters to a group of size  $2m - 1$ .
- . Without additional voters, a majority of  $2m - 1$  (i.e. at least  $m$  voters) makes the right decision with probability  $M_{2m-1}(p)$ .
- . Note that  $M_{2m-1}(p) = \sum_m^{2m-1} b_{2m-1}(x; p)$ .
- . For  $2m + 1$  voters, there is a majority for the right decision,
  - . if at least  $m + 1$  of  $2m - 1$  voters vote correctly,
  - . or if exactly  $m$  of  $2m - 1$  voters vote correctly and at least 1 of the 2 other voters votes correctly,
  - . or if exactly  $m - 1$  of  $2m - 1$  voters vote correctly and both the 2 other voters vote correctly.

. Thus, we have:

$$M_{2m+1}(p) = M_{2m-1}(p) - b_{2m-1}(m; p) \\ + b_{2m-1}(m; p)(p^2 + 2p(1 - p)) + b_{2m-1}(m - 1; p)p^2.$$

. Let  $q \equiv 1 - p$ . Since  $b_{2m-1}(m - 1; p) = b_{2m-1}(m; p)(q/p)$ , we can simplify the above expression and get:

$$M_{2m+1}(p) = M_{2m-1}(p) + q(p - q)b_{2m-1}(m; p). \quad (1)$$

**Condorcet Jury Theorem** For  $p > (<)1/2$ , the majority function  $M_{2m+1}(p)$  is monotonically increasing (resp. decreasing) in  $m$  and  $\lim_{m \rightarrow \infty} M_{2m+1}(p) = 1$  ( $= 0$ ).  $M_{2m+1}(1/2) = 1/2$ , for all  $m$ . For  $p \in (1/2, 1)$ , we have:  $M_{2m+1}(p) > p$ .

- . If  $n = 2m + 1$  voters act independently and each one decides correctly with probability  $p > 1/2$ , then the probability for majority voting to get the right decision converges to 1 for  $n \rightarrow \infty$ .
- . The convergence is fast, e.g. it is  $> 0,99$  for  $p = 0.8$  and  $n = 13$ .
- . We have “vox populi, vox dei”, i.e. majority voting will be almost never wrong if the number of voters is sufficiently large.
- . It may be reasonable that a decision is made by a group even if every member of the group has lower competence (but  $p > 1/2$ ) than a single “competent expert” who would decide alone.

- . The proof is made by means of the recursion formula (1):

$$M_{2m+1}(p) = M_{2m-1}(p) + q(p - q)b_{2m-1}(m; p)$$

- . For  $p > 1/2$ ,  $M_{2m+1}(p)$  increases monotonically in  $m$ , as  $q(p - q)$  is positive.
- . For  $p < 1/2$ ,  $M_{2m+1}(p)$  decreases monotonically in  $m$ , as  $q(p - q)$  is negative.
- . For  $p = 1/2$ ,  $M_{2m+1}(p)$  is constant in  $m$ , as  $q(p - q) = 0$ .
- . The proof for  $\lim_{m \rightarrow \infty} M_{2m+1}(p) = 1$  is made by expansion of the recursion formula. The calculation is very extensive.
- . The property  $\lim_{m \rightarrow \infty} M_{2m+1}(p) = 1$  can also be derived directly by the law of large numbers.

## Extensions

- . Is there really a “right” and a “wrong” decision?

Voters may have different preferences.

- . Voters differ in their competence and information.

Each voter  $i$ 's signal's precision may take a different value  $p_i$ .

- . Voters observe common information: signals may be correlated.

- . The model assumes that voters vote “truthfully.”

- . Each voter  $i$  observes a signal  $s_i \in \{0, 1\}$ , which is “correct” with probability  $p$ . Then  $i$  votes  $v_i = s_i$ .

- . Is this behavior compatible with equilibrium?

## Strategic Voting (Austen-Smith and Banks, 1996)

- . 3 voters choose an alternative  $d \in \{0, 1\}$  by majority.
- . Given unknown state  $x \in \{0, 1\}$ , each voter  $j$ 's payoff is

$$u_j(x, d) = \begin{cases} 0 & \text{if } d = x \\ -1 & \text{if } d \neq x. \end{cases}$$

- . The prior is favorable to the status quo:  $\Pr(x = 0) = \pi > 1/2$ .
- . Each  $j$  has a private signal  $s_j \in \{0, 1\}$ ,  $\Pr(s_j = x|x) = p > 1/2$ , signals are independent across voter.
- . Each  $j$ 's voting strategy is a function  $v_j : \{0, 1\} \rightarrow \{0, 1\}$  that maps each signal into a vote 0 or 1.
- . I show that if  $\pi$  is large relative to  $p$ , then truthful voting is not a Bayesian equilibrium.

- . Consider a voter  $j$  with signal  $s_j = 1$ . Suppose by contradiction that the other voters  $k$  and  $\ell$  vote truthfully.
- . Let  $r_j$  be probability of  $x = 1$ , given  $j$ 's equilibrium information.
- . Her expected payoff is  $-r_j$  for  $d = 0$  and  $-(1 - r_j)$  for  $d = 1$ . Voter  $j$  prefers  $d = 0$  if  $r_j < 1/2$ , and  $d = 1$  if  $r_j > 1/2$ .
- .  $j$ 's vote has no effect on  $d$ , unless one other voter votes  $v_k = 0$  and the other one votes  $v_\ell = 1$ .
- . Hence, voter  $j$  with  $s_j = 1$  votes  $v_j = 1$  if and only if

$$r_j = Pr(x = 1 | s_j = s_\ell = 1, s_k = 0) = \frac{(1-p)p^2(1-\pi)}{(1-p)p^2(1-\pi) + p(1-p)^2\pi} > \frac{1}{2}.$$

- . If  $\pi$  is large relative to  $p$ , then  $j$  votes  $v_j = 0$  although  $s_j = 1$ .
- . This happens even if truthful voting would be Pareto superior,

$$r_j = Pr(x = 1 | s_j = s_k = s_\ell = 1) = \frac{p^3(1-\pi)}{p^3(1-\pi) + (1-p)^3\pi} > \frac{1}{2}.$$

## Condorcet Jury Theorem (Feddersen and Pesendorfer, 1997)

### **The model**

- .  $n + 1$  voters must make a decision  $d = 0, 1$ .
- . Each voter  $i$ 's payoff is  $u(d, x, b_i)$ ,
- $x \in X = [0, 1]$  is an unknown state with full support density  $g$ ,
- $b_i \in B = [-1, 1]$  is private information bias, full support density  $f$ .
- . Let  $v(x, b) \equiv u(1, x, b) - u(0, x, b)$  be the utility difference,  
 $v(x, b)$  is continuous and strictly increasing,  
 $v(-1, b) < 0$  and  $v(1, b) > 0$  for all  $b$ .

- . Each voter receives a signal  $s \in \{\underline{s}, \dots, \bar{s}\} \equiv S$ , with full support probability  $p(s|x)$ , continuous in  $s$  for all  $x$ .
- . Monotone Likelihood Ratio Property: If  $s > s'$  and  $x > x'$ , then  $p(s'|x')p(s|x) > p(s|x')p(s'|x)$ .
- . Given quorum  $q \in [1/2, 1)$ , each voter  $i$  votes  $v_i = 0, 1$ .
- . The voting outcome is  $d = 1$  iff  $\#\{i : v_i = 1\} \geq (n + 1)q$ .
- . A mixed strategy for voter  $i$  is  $\sigma_i : B \times S \rightarrow [0, 1]$ .
- . We consider weakly undominated symmetric Nash equilibria.

## Analysis

- . A voter  $i$ 's vote influences the election outcome iff one vote is pivotal: exactly  $qn$  of the other  $n$  voters voted  $v = 1$ .
- . The “average” probability that a voter votes 1 in state  $x$  is

$$\tau(x, \sigma) = \sum_{s=\underline{s}}^{\bar{s}} p(s|x) \int_B \sigma(b, s) f(b) db.$$

- . The probability that a vote is pivotal in state  $x$  is:

$$\Pr(\text{piv}|x, \sigma) = \binom{n}{qn} \tau(x, \sigma)^{qn} (1 - \tau(x, \sigma))^{n-qn}.$$

- . When  $0 < \tau(x, \sigma) < 1$ , we have  $\Pr(\text{piv}|x, \sigma) > 0$  for all  $x$ .

. Densities of state  $x$  conditional on event  $piv$ , and on signal  $s$  are:

$$g(x|piv, \sigma) = \frac{\Pr(piv|x, \sigma)g(x)}{\int_X \Pr(piv|x, \sigma)g(x) dx},$$
$$g(x|s, piv, \sigma) = \frac{p(s|x)g(x|piv, \sigma)}{\int_X p(s|x)g(x|piv, \sigma) dx}.$$

. As  $s$  satisfies MLRP,  $g(x|s, piv, \sigma)$  is first-order stochastically increasing in  $s$ , and  $E[v(b, x)|s, piv, \sigma]$  increases in  $s$ .

. Hence, every voting equilibrium  $\sigma$  is characterized by ordered cutpoints  $(b_s)_{s \in S}$  such that  $-1 < b_{\bar{s}} < \dots < b_{\underline{s}} < 1$  and

$$E[v(b_s, x)|s, piv, \sigma] = 0 \text{ for all } s.$$

. For all  $s$ ,  $\sigma(b, s) = 0$  if  $b < b_s$  and  $\sigma(b, s) = 1$  if  $b > b_s$ , and  $0 < \tau(x, \sigma) < 1$  increases in  $x$ . The election is informative.

- . Suppose that the number of voters grows to infinity.
- . Then, the expected fraction of voters who vote informatively in equilibrium must converge to zero, and the election must be close.

**Theorem 1** Let  $(\sigma^n)_{n \geq 1}$  be a sequence of voting equilibria, and let  $((b_s^n)_{s \in S})_{n \geq 1}$  be the corresponding cutpoints. Then  $b_{\underline{s}}^n - b_{\bar{s}}^n \rightarrow 0$ .

*Sketch of Proof:* In equilibrium, each voter chooses as if she was pivotal, i.e., as if  $qn$  out of  $n$  voters voted  $v = 1$ .

- . Equilibrium beliefs about  $\tau(x, \sigma)$  must be concentrated around  $q$ .
- . Beliefs about the state  $x$  must concentrate on states  $x'$  such that  $\tau(x', \sigma)$  is close to  $q$ , regardless of what the true state  $x$  is.
- . Regardless of the state, the election must be close.
- . If the fraction of voters who voted informatively did not vanish, then the election would not be close for all states.

- . Large elections almost always choose the alternative that would have been chosen if the state  $x$  were common knowledge.
- . Let  $b^* = F^{-1}(q)$  be expected bias of the “pivotal” voter.
- . Let  $x^*$  be the marginal state such that  $v(b^*, x) = 0$ .

**Theorem** Every sequence of voting equilibria  $(\sigma^n)_{n \geq 1}$  is such that  $\Pr(x < x^*, d^n = 1) \rightarrow 0$  and  $\Pr(x > x^*, d^n = 0) \rightarrow 0$ .

The probability of a decision contrary to the pivotal voter's preference vanishes as  $n$  grows to infinity.

*Sketch of Proof:* Conditional on pivotal voting, the distribution over states puts almost all weight close to one state  $x^n$ .

- . If  $v(b^*, x^n) < \varepsilon < 0$ , then the fraction of votes for 1 in state  $x^n$  would be boundedly smaller than  $q$ : election would not be close.
- . We conclude that  $x^n \rightarrow x^*$  as  $n \rightarrow \infty$ .

## Swing voter's curse (Feddersen and Pesendorfer, 1996)

- . Elections aggregate individual preferences and information.
- . Information of common value, but some voters are not informed.
- . Uninformed voters abstain, to avoid swinging the election against common interest.
- . In fact, many voters do not vote, although the cost of voting is often negligible.
- . Strategic abstention delivers first best.
- . The winning candidate is the same as if all voters knew all voters' information.

## The model

- . There are 2 states  $\omega = 0, 1$ , with  $\pi = \Pr(\omega = 0) \geq 1/2$ , and 2 party candidates  $j = 0, 1$ , with platforms  $x_j = 0, 1$ .
- . There are  $N + 1$  possible voters, each votes with prob.  $1 - p_A$ .
- . With prob.  $p_0$  (prob.  $p_1$ ), a voter is partisan for party 0 (party 1).
- . With probability  $p_n = 1 - p_0 - p_1$  the voter is independent: her utility is  $u_n(x, \omega) = -|x - \omega|$ .
- . Each voter receives a signal  $s \in S = \{0, a, 1\}$ .
- . With probability  $1 - q$ ,  $s$  is uninformative and equal to  $a$ .
- . When signal  $s$  is informative,  $\Pr(s = \omega | \omega) = p > 1/2$ .
- . Each voter chooses  $v \in \{0, A, 1\}$ , where  $A$  is abstention.

- . I focus on symmetric Nash equilibria: voters with same type and signal vote the same candidate.
- . In equilibrium, type-0 (type-1) voters vote  $v_0 = 0$  ( $v_1 = 1$ ).
- . All informed independents vote according to their signal:  
 $v_n(s) = s$  if  $s = 0, 1$ .
- . The mixed strategy of uninformed independent agents (UIAs) is  $\sigma = (\sigma_0, \sigma_1, \sigma_A) \in \Delta^3$ .

## Equilibrium

. Given the strategy  $\sigma$ , let  $\rho_{\omega,j}(\sigma)$  be the probability of a vote for  $j$  if the state is  $\omega$  is as follows

$$\rho_{\omega,j}(\sigma) = p_j + p_n(1 - q)\sigma_j + p_nq(1 - p) \quad \text{if } \omega \neq x_j,$$

$$\rho_{\omega,j}(\sigma) = p_j + p_n(1 - q)\sigma_j + p_nqp \quad \text{if } \omega = x_j.$$

. Let  $\rho_{\omega,A}(\sigma)$  be the probability of an abstention if the state is  $\omega$ :

$$\rho_{0,A}(\sigma) = \rho_{1,A}(\sigma) = \rho_A(\sigma) = p_n(1 - q)\sigma_A + p_A.$$

. For any voter, the probability of a tie among the other voters is:

$$P_T^{\omega,\sigma} = \sum_{\ell=0}^{N/2} \frac{N!}{\ell!\ell!(N-2\ell)!} \rho_{\omega,A}(\sigma)^{N-2\ell} \rho_{\omega,0}(\sigma)^\ell \rho_{\omega,1}(\sigma)^\ell.$$

. The probability that candidate  $j$  is down by 1 vote is:

$$P_j^{\omega,\sigma} = \sum_{\ell=0}^{(N/2)-1} \frac{N!}{(\ell+1)!\ell!(N-2\ell-1)!} \rho_{\omega,A}(\sigma)^{N-2\ell-1} \rho_{\omega,1-j}(\sigma)^{\ell+1} \rho_{\omega,j}(\sigma)^\ell.$$

. Let  $Eu_n(v, \sigma)$  be an UIA expected payoff of voting  $v$ , when the other voters use  $\sigma$ :

$$Eu_n(1, \sigma) - Eu_n(A, \sigma) = \frac{1}{2}[(1 - \pi)(P_T^{1,\sigma} + P_1^{1,\sigma}) - \pi(P_T^{0,\sigma} + P_1^{0,\sigma})]$$

$$Eu_n(0, \sigma) - Eu_n(A, \sigma) = \frac{1}{2}[\pi(P_T^{0,\sigma} + P_0^{0,\sigma}) - (1 - \pi)(P_T^{1,\sigma} + P_0^{1,\sigma})].$$

$$Eu_n(1, \sigma) - Eu_n(0, \sigma) = (1 - \pi)[P_T^{1,\sigma} + \frac{1}{2}(P_1^{1,\sigma} + P_0^{1,\sigma})] \\ - \pi[P_T^{0,\sigma} + \frac{1}{2}(P_1^{0,\sigma} + P_0^{0,\sigma})].$$

**Proposition** Suppose  $p_A > 0$ ,  $q > 0$ ,  $N \geq 2$  and  $N$  even.  
For any symmetric  $\sigma$  s.t no voter plays a strictly dominated strategy,  $Eu_n(1, \sigma) = Eu_n(0, \sigma)$  implies  $Eu_n(1, \sigma) < Eu_n(A, \sigma)$ .

. An UIA strictly prefers to abstain whenever indifferent between voting for 1 or 0, and no voter uses a strictly dominated strategy.

. This is the swing voter's curse.

. To consider large elections, define a sequence of games with  $N + 1$  voters and associated strategy profiles  $\{\sigma^N\}_{N=0}^{\infty}$ .

**Proposition** Suppose  $q > 0$ ,  $p_n(1 - q) < |p_0 - p_1|$  and  $p_A > 0$ . Let  $\{\sigma^N\}_{N=0}^{\infty}$  be a sequence of equilibria.

. If  $p_n(1 - q) < p_0 - p_1$  then  $\lim_{N \rightarrow \infty} \sigma_1^N = 1$ , i.e., all UIAs vote for candidate 1.

. If  $p_n(1 - q) < p_1 - p_0$  then  $\lim_{N \rightarrow \infty} \sigma_0^N = 1$ , i.e., all UIAs vote for candidate 0.

. The swing voter's curse can lead to large scale abstention by the UIAs in large elections.

. This happens when the expected fraction of UIAs is too small to compensate for a candidate partisan advantage.

. Instead, when the fraction of UIAs is large enough to offset partisan bias, there are no pure strategy equilibria.

- . UIAs mix between abstention and voting against the difference in partisan support to compensate exactly.
- . The equilibrium winning candidate is approximately the same as the candidate that would win if voters had perfect information.

**Proposition** Suppose  $q > 0$ ,  $p_n(1 - q) \geq |p_0 - p_1|$  and  $p_A > 0$ . Let  $\{\sigma^N\}_{N=0}^\infty$  be a sequence of equilibria.

- . If  $p_n(1 - q) \geq p_0 - p_1 > 0$  then UIAs mix between voting for candidate 1 and abstaining, with  $\lim \sigma_1^N = \frac{p_0 - p_1}{p_n(1 - q)}$ .
- . If  $p_n(1 - q) \geq p_1 - p_0 > 0$  then UIAs mix between voting for candidate 0 and abstaining, with  $\lim \sigma_1^N = \frac{p_1 - p_0}{p_n(1 - q)}$ .
- . If  $p_0 - p_1 = 0$  then UIAs abstain:  $\lim \sigma_A^N = 1$ .
- . For every  $\epsilon$  there exists an  $N$  such that for  $\bar{N} > N$  the probability that equilibrium fully aggregates information is greater than  $1 - \epsilon$ .

## Summary

- . I have considered how well elections aggregate information.
- . If voters vote truthfully, then they select the “best” alternative by the law of large numbers.
- . The fraction of voters who vote informatively in equilibrium converges to zero in large elections, and the election must be close.
- . Nevertheless the chosen alternative is the same that would be chosen if all information became common knowledge.
- . I have presented a model in which voters have different information about candidates’ valence.
- . There exists an equilibrium in which informed non-partisan voters are pivotal, and the “best” candidate is elected.

## Next lecture

- . We will review models of cheap talk and political advice.
- . Congress may benefit from committing not to amend a committee's bill proposal, and put it to vote against the status quo.
- . Unless the status quo is in line with the committee's bias, it disciplines the committee's proposal. (Gilligan and Krehbiel 1987).
- . If an expert's loyalty is uncertain, repeated information transmission yields reputational concerns.
- . Reputational concerns may lead to more disclosure but also to "political correctness" and conformism (Morris 2001).
- . When information is verifiable, beliefs divergent from the DM act as incentives for information acquisition (Che and Kartik 2009).
- . This incentive is reinforced by preference divergences, and dominates information withholding unless beliefs diverge too much.